

Failure Probabilities in Computing and Data Transmission

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SA, Fiorini, Lipsky, Rolski, Sheahan

Asymptotic total times for task that must restart after
a failure occurs

Mathematics of Operations Research 2008/9

SA, Lars N. Andersen

Failure recovery, parallel computing, and extreme values

Journal of Statistical Methodology and Applications

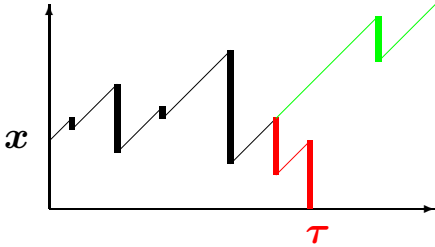
2 working papers (checkpointing, importance sampling)

Related work: [Jelenkovic & Tan](#)

Cramér-Lundberg Theory

$$\psi(x) \sim C e^{-\gamma x}$$

CL: transforms, complex analysis



Classical expression P-K-B-B

Pollaczek-Khintchine-Beekman-Bowers

β Poisson int'y, unit premium rate

B claim distr'n, integrated tail B_I :

$$B_I(x) = \frac{1}{\mu} \int_0^x \bar{B}(y) dy$$

$$\psi(u) = (1 - \rho) \sum_{n=0}^{\infty} \rho^n \bar{B}_I^{*n}(u)$$

$$\rho = \beta \mu_B < 1$$

Geometric sum

Geometric Sums

$$S_N = V_1 + \cdots + V_N$$

$$V_n \sim F, N \text{ geom}(\rho)$$

$$Z(x) = P(S_N > x) \text{ solves}$$

$$Z(x) = \rho \bar{F}(x) + \int_0^x Z(x-y) \rho F(dy) \quad (1)$$

Defective renewal equation

Solution exponentially decaying

$\rho F(dy)$ has mass $\rho < 1$

Choose γ as solution of

$$\int_0^\infty e^{\gamma x} \rho F(dx) = 1$$

Let

$$F^*(dx) = e^{\gamma x} \rho F(dx)$$

$$Z^*(x) = e^{\gamma x} Z(x)$$

$$z^*(x) = e^{\gamma x} \rho \bar{F}(x)$$

Multiply (1) by $e^{\gamma x}$

$$Z(x) \sim C e^{-\gamma x}$$

RESTART

T ideal job time $\sim F$
(program time; file length;
call center time)

U failure time $\sim G$
(often exponential)

X total time $\sim H$

$$\frac{\frac{\frac{T}{U_1}}{U_2}}{U_N} \quad X = T + U_1 + \dots + U_N$$
$$\frac{\quad}{U_{N+1}}$$

Target: tail $\overline{H}(x) = P(X > x)$

T bounded $\Rightarrow \overline{H}(x) \approx e^{-\gamma x}$

T unbd $\Rightarrow H$ heavy-tailed

$$T \equiv t$$

$$\begin{array}{c}
 \hline
 t \\
 \hline
 U_1 \\
 \hline
 U_2 \\
 \hline
 U_N \\
 \hline
 U_{N+1} \\
 \hline
 \end{array}
 \quad D = t + U_1 + \dots + U_N$$

Geometric sum

Success probability $\bar{G}(t)$

Summands: $U \mid U \leq t$

$$\gamma(t) \text{ solution of } 1 = \int_0^t e^{\gamma(t)u} g(u) du$$

$$P(S > x) \sim C(t) e^{-\gamma(t)x}$$

$$\bar{H}(x) \sim e^{\gamma(t)t} C(t) e^{-\gamma(t)x}$$

$$\gamma(t) \text{ solution of } 1 = \int_0^t e^{\gamma(t)u} g(u) du$$

$$\gamma(t) \sim \mu \bar{G}(t) \quad \mu = 1/EU$$

Crucial lemma:

$$\begin{aligned} \bar{H}(x) &= \int_0^\infty \bar{H}_t(x) f(t) dt \\ &\sim \int_{t_0}^\infty \exp\{-\mu \bar{G}(t)x\} f(t) dt \end{aligned}$$

Purely analytical problem;
non-trivial

Diagonal case $F = G$:

$$\begin{aligned} y = \mu G(t)x &\Rightarrow dy = \mu x f(t) dt \\ \frac{1}{\mu x} (1 - \exp\{-\mu \bar{G}(t_0)x\}) &\sim \frac{1}{\mu x} \end{aligned}$$

General idea:

same substitution

assume F, G are connected s.t.

this leads to integral of

Abelian/Tauberian type.

$$f(t) = g(t)\overline{G}(t)^{\beta-1}L_0(\overline{G}(t))$$

$$\overline{H}(x) \sim \frac{\Gamma(\beta)L_0(1/x)}{\mu^\beta x^\beta}$$

f, g regularly varying
of form $L(t)/t^{1+\alpha}$

$\alpha_F, \alpha_G, L_F, L_G$

$$\overline{H}(x) = L_H(x)/x^{\alpha_H}$$

$$\alpha_H = \alpha_F/\alpha_G, L_H(x) \sim c_H L'_H(x)$$

$$c_H = \frac{\Gamma(\alpha_H)(\alpha_G + 1)^{\alpha_H-1}}{\mu^{\alpha_H}}$$

$$L'_H(x) = \frac{L_F(x^{1/\alpha_G})}{L_G^{\alpha_H}(x^{1/\alpha_G})}.$$

f, g

of form $\exp\{-\lambda t^\eta\} t^{1+\alpha} L(t)$

$\lambda_F, \alpha_F, \lambda_G, \alpha_G, L_F, L_G$

$$\overline{H}(x) = \log^{\kappa_H} x L_H(x) / x^{\alpha_H}$$

$$\alpha_H = \lambda_F / \lambda_G, L_H(x) \sim c_H L'_H(x)$$

$$c_H = \frac{\Gamma(\alpha_H)}{\mu^{\alpha_H} \lambda_G^{\omega_H + \alpha_H - 1}}$$

$$\omega_H = \alpha_F / \eta - \alpha_G \lambda_F / \lambda_G \eta$$

$$\kappa_H = \omega_H + (1 - 1/\eta)(\lambda_F / \lambda_G - 1)$$

$$L'_H(x) = \frac{L_F(\log x^{1/\eta})}{L_G^{\alpha_H}(\log x^{1/\eta})}.$$

F Gamma-like: $\overline{F}(x) \sim Ax^\eta e^{-\delta x}$

$$g(t) = \beta e^{-\beta t}$$

$$\overline{H}(x) \sim \frac{A\Gamma(\delta/\beta) \log^\eta x}{\beta^{\delta/\beta - 1 - \eta} x^{\delta/\beta}}$$

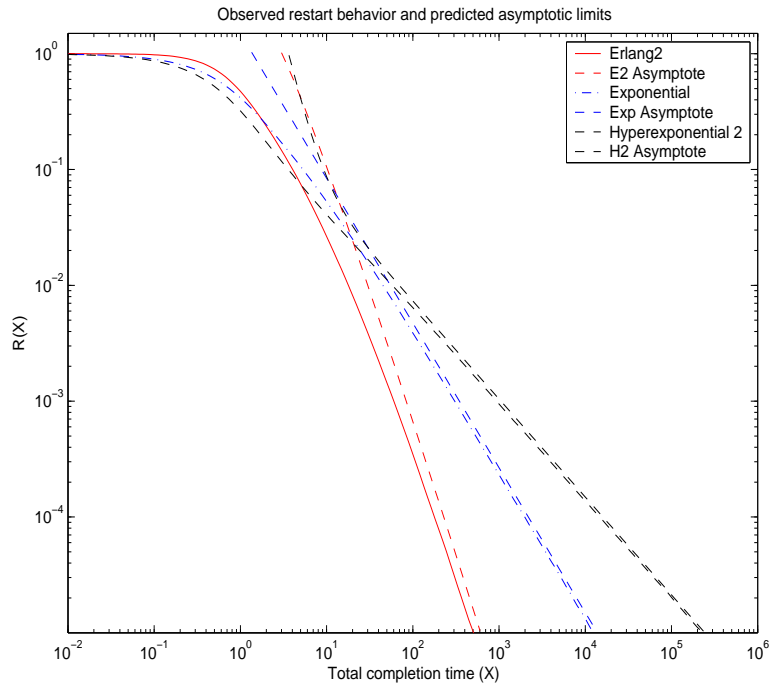


Figure 1:

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Crude Monte Carlo
Computationally heavy
Importance sampling

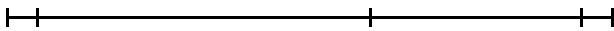
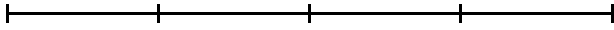
4 examples of each of F, G :
 LT Weibull
 exponential
 HT Weibull
 power

$\cdot \bar{F}(t)$ $\bar{G}(u)$	e^{-t^2}	e^{-t}	$e^{-t^{1/2}}$	$\frac{1}{t^\alpha}$
e^{-u^2}	$\frac{1}{x}$	$e^{-\log^{1/2} x}$	$e^{-\log^{1/4} x}$	$\frac{1}{\log^{\alpha/2} x}$
e^{-u}	$e^{-\log^2 x}$	$\frac{1}{x}$	$e^{-\log^{1/2} x}$	$\frac{1}{\log^\alpha x}$
$e^{-u^{1/2}}$	$e^{-\log^4 x}$	$e^{-\log^2 x}$	$\frac{1}{x}$	$\frac{1}{\log^{2\alpha} x}$
$\frac{1}{u^\alpha}$	$e^{-x^{\frac{2}{2+\alpha}}}$	$e^{-x^{\frac{1}{1+\alpha}}}$	$e^{-x^{\frac{1/2}{1/2+\alpha}}}$	$\frac{1}{x}$

Constants omitted $e^{-c \log^{1/2} x}$; $\frac{1}{x} = e^{-\log x}$
 In some corners even $\log \log$ asymptotics

Fragmentation

K parts



Parallel computing:

$$X = \max(X_1, \dots, X_K)$$

Checkpointing:

$$X = X_1 + \dots + X_K$$

Parallel computing

T_1, \dots, T_K i.i.d. Gamma(α_K, λ)

Exponential(μ) failures

$\alpha_K \equiv K$: $\mathbb{P}(X_i > x) \sim c \log^{\alpha-1} / x^r$,

$r = \lambda/\mu$

$\frac{L(K)}{K^\beta} X \rightarrow \text{Frechet}$

$\alpha_K \rightarrow \infty$:

$X \approx e^{r\alpha_K}$

$\frac{\log X}{r\alpha_K} \rightarrow c$

Checkpointing

K segments of lengths h_1, \dots, h_K

Checkpoints $t_0 = 0, t_1 = h_1, t_2 = h_1 + h_2, \dots$

A: T is deterministic, $T \equiv t$, and the checkpoints are deterministic and equally spaced, $t_1 = t/K, t_2 = 2t/K, \dots, t_{K-1} = (K - 1)t/K$. Equivalently, $h_k = t/K$.

B: T is deterministic, $T \equiv t$, and the checkpoints are deterministic but not equally spaced, for simplicity $h_k \neq h_\ell$ for $k \neq \ell$.

C: T is deterministic, $T \equiv t$, and the checkpoints are random. More precisely, the set $\{t_1, \dots, t_{K-1}\}$ is the outcome of $K - 1$ i.i.d. uniform r.v.'s on $(0, t)$. That is, $t_1 < \dots < t_{K-1}$ are the order statistics of $K - 1$ i.i.d. uniform r.v.'s on $(0, t)$.

D: T is random and the checkpoints equally spaced, $h_k \equiv h$. Thus, $K = \lceil T/h \rceil$ is random

E: T is random and the checkpoints are given by $t_k = t'_k T$ for a deterministic set of constants $0 = t'_0 < t'_1 < \dots < t'_{K-1} < 1$.

Model A:

$$\mathbb{P}(X > x) \sim c_K x^{K-1} e^{-x\gamma(t/K)}$$

Let $U_1, \dots, U_m \geq 0$ be independent r.v.'s such that $\mathbb{P}(U_i > x) \sim C_i x^{\alpha_i-1} e^{-\eta x}$ for some $\eta > 0$, $\alpha_1, \dots, \alpha_m > 0$ and some C_1, \dots, C_m . Then

$$\mathbb{P}(U_1 + \dots + U_m > x) \sim C x^{\alpha-1} e^{-\eta x},$$

where $\alpha = \alpha_1 + \dots + \alpha_m$, $C = C_1 \Gamma(\alpha_1) \dots C_m \Gamma(\alpha_m) / \Gamma(\alpha)$

Model B:

$$\mathbb{P}(X > x) \sim c e^{-x\gamma(h^*)}$$

$$h^* = \max(h_1, \dots, h_K)$$

Let $U_1, U_2 \geq 0$ be independent r.v.'s such that $\mathbb{P}(U_1 > x) \sim C e^{-\eta x}$ and $\mathbb{P}(U_2 > x) = o(e^{-\eta x})$ for some $\eta > 0$. Then

$$\mathbb{P}(U_1 + U_2 > x) \sim C \mathbb{E} e^{\eta U_2} e^{-\eta x}.$$

Model C

Assume that the checkpoints are $0, t$ and $K - 1$ i.i.d. uniform r.v.'s on $(0, t)$. That is, the checkpoints are $0 = t_0 < t_1 < \dots < t_{K-1} < t_K = t$ where $t_1 < \dots < t_{K-1}$ are the order statistics of $K - 1$ i.i.d. uniform r.v.'s R_1, \dots, R_{K-1} on $(0, t)$. Then

$$\mathbb{P}(X^* > x) \sim C_7(t) \frac{e^{-\gamma(t)x}}{x^{K-1}}$$

$$\begin{aligned} \mathbb{P}(h_{k^*} > y) &= (K - 1)(1 - y)^{K-2} \\ &\quad - \binom{K - 1}{2} (1 - 2y)^{K-2} \\ &\quad + \dots + (-1)^{i-1} \binom{K - 1}{i} (1 - iy)^{K-2} \end{aligned}$$

Fisher 1929

$$\begin{aligned} \mathbb{P}(h_{k^*} \in dy) &= (K - 1)\mathbb{P}(R_1 \in dy, k^* = 1) \\ &= \mathbb{P}(R_1 \in dy, R_2 > y, \dots, R_{K-1} > y) \\ &= \frac{K - 1}{t^{K-1}} (1 - y)^{K-2}. \end{aligned}$$

Model D

Assume that the distribution of T is regularly varying, $\mathbb{P}(T > t) = L(t)/x^\alpha$ with $\alpha > 0$ and $L(\cdot)$ slowly varying. Then

$$\begin{aligned}\mathbb{P}(X > x) &\sim \mathbb{P}(T > xh/m_1(h)) \\ &\sim \frac{m_1(h)^\alpha L(x)}{h^\alpha x^\alpha} \quad (2)\end{aligned}$$

Asmussen, Klüppelberg & Sigman 1999

Assume $f(t) \sim ct^\alpha e^{-\lambda t}$ as $t \rightarrow \infty$ where $-\infty < \alpha < \infty$. Then $\mathbb{P}(X > x) \sim C_7 e^{-\gamma_2 x}$ where $\gamma_2 = \gamma_2(h)$ is the solution of $\widehat{H}[\gamma_2] = e^{\lambda h}$ and $C_7 = \dots$

Let $U_1, U_2, \dots > 0$ be i.i.d. and define $S_n = U_1 + \dots + U_n$. Let $N \in \mathbb{N}$ be an independent r.v. such that $\mathbb{P}(N \geq n) \sim cn^{\alpha-1} \rho^n$ with $\alpha > 0$ and $0 < \rho < 1$. Assume that a solution $\theta > 0$ of $\rho \mathbb{E} e^{\theta U} = 1$ exists with the additional property $\mathbb{E} e^{\theta' U} < \infty$ for some $\theta' > \theta$. Then $\mathbb{P}(S_N > x) \sim C_9 x^{\alpha-1} e^{-\theta x}$, where $\theta > 0$ is the solution of $\rho \mathbb{E} e^{\theta U} = 1$

and

$$C_9 = \frac{c(1 - \rho)}{\theta m_1(\theta)^\alpha} \quad \text{where} \quad m_1(\theta) = \rho \mathbb{E}[U e^\theta]$$