

# Lecture 9 Decomposition and Distributed Optimization

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## Aims

After this lecture, you should be able to

- bound optimal value by solving (convex but maybe non-smooth) dual
- compute subgradients and use the subgradient method
- use dual decomposition for efficient large-scale optimization
- derive distributed optimization schemes using decomposition techniques

## Disposition

- Lagrange duality review
- Subgradients and non-differentiable optimization
- Dual optimization
- Distributed optimization
- Primal decomposition

## The Lagrangian

Optimization problem in standard form

$$\begin{aligned} &\text{minimize} && f_0(x) \\ &\text{subject to} && f_i(x) \leq 0 \quad i = 1, \dots, m \\ & && h_i(x) = 0 \quad i = 1, \dots, p \end{aligned}$$

variable  $x \in \mathbb{R}^n$ , domain  $\mathcal{D}$ , optimal value  $p^*$ .

*Lagrangian*:  $L : \mathbb{R}^n \times \mathbb{R}^m \times \mathbb{R}^p \mapsto \mathbb{R}$  with domain  $\mathcal{D} \times \mathbb{R}^m \times \mathbb{R}^p$

$$L(x, \lambda, \nu) = f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

Weighted sum of objective and constraint functions,

- $\lambda_i$  is *Lagrange multiplier* (or *dual variable*) associated with  $f_i(x) \leq 0$
- $\nu_i$  is *Lagrange multiplier* associated with  $h_i(x) = 0$

## Lagrange dual: concave and lower bounding

*Lagrange dual function*:  $g : \mathbb{R}^m \times \mathbb{R}^p \mapsto \mathbb{R}$ ,

$$\begin{aligned} g(\lambda, \nu) &= \inf_{x \in \mathcal{D}} L(x, \lambda, \nu) \\ &= \inf_{x \in \mathcal{D}} \left( f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) + \sum_{i=1}^p \nu_i h_i(x) \right) \end{aligned}$$

**Fact:**  $g$  is concave (can be  $-\infty$  for some  $\lambda, \nu$ )

**Fact (lower bound property):** if  $\lambda \geq 0$ , then  $g(\lambda, \nu) \leq p^*$

## Dual problem: best lower bound

*Lagrange dual problem*:

$$\begin{aligned} &\text{maximize} && g(\lambda, \nu) \\ &\text{subject to} && \lambda \geq 0 \end{aligned}$$

Finds *best lower bound*  $d^*$  on  $p^*$  obtained from Lagrange dual function

Observations:

- a convex optimization problem; optimal value denoted  $d^*$
- $\lambda, \nu$  are dual feasible if  $\lambda \geq 0$  and  $(\lambda, \nu) \in \text{dom } g$
- often simplified by making the constraint  $(\lambda, \nu) \in \text{dom } g$  explicit

## Duality: structure of solution

If strong duality holds, optimal  $(x^*, \lambda^*, v^*)$  must satisfy KKT conditions:

1. Primal constraints:

$$f_i(x^*) \leq 0, \quad i = 1, \dots, m, \quad h_i(x^*) = 0, \quad i = 1, \dots, p$$

2. Dual constraints:

$$\lambda^* \geq 0$$

3. Complementary slackness:

$$\lambda_i^* f_i(x^*) = 0, \quad i = 1, \dots, m$$

4. Gradient of Lagrangian w.r.t.  $x$  vanishes

$$\nabla f_0(x^*) + \sum_{i=1}^m \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^p \nu_i^* \nabla h_i(x^*) = 0$$

## Duality: sensitivity analysis

Nominal and perturbed problem

$$\begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq 0 \quad i = 1, \dots, m \end{array} \quad \begin{array}{ll} \text{minimize} & f_0(x) \\ \text{subject to} & f_i(x) \leq u_i \quad i = 1, \dots, m \end{array}$$

$p^*(u)$  is optimal value as function of  $u$

What can we say about  $p^*(u)$  by solving the nominal primal and its dual?

By weak duality,

$$p^*(u) \geq p^*(0) - \sum_{i=1}^m \lambda_i^* u_i$$

If  $p^*(u)$  differentiable

$$\lambda_i = -\frac{\partial p^*(0)}{\partial u_i}$$

$\lambda_i$  is the *sensitivity* of  $p^*$  with respect to  $i$ th constraint.

## Duality: efficient optimization

Today's lecture: duality sometimes enables efficient optimization

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## Subgradients

A *subgradient*  $h$  of a convex function  $f$  at  $x$  is any vector that satisfies

$$f(y) \geq f(x) + h^T(x)(y - x) \quad \text{for all } y$$

Subgradients

- gives affine global underestimator of  $f$
- if  $f$  is convex and differentiable,  $\nabla f(x)$  is a subgradient of  $f$  at  $x$

The set of subgradients at  $x$  is called the subdifferential, denoted  $\partial f(x)$

**Quiz:** determine  $\partial f(x)$  for  $f(x) = |x|$ .

## Subgradient method

For convex function  $f$ , subgradient method uses iteration

$$x(t+1) = x(t) - \alpha(t)h(x(t))$$

where

- $x(t)$  is  $t$ th iterate,  $g^{(k)}$  any subgradient,  $\alpha(t) > 0$  is step-length

Guaranteed to converge if (for example)  $h$  bounded and  $\alpha$  satisfies

$$\sum_t \alpha^2(t) < \infty, \quad \sum_{t=1}^{\infty} \alpha(t) = \infty$$

**Proof:**

Common choice:

$$\alpha(t) = \frac{a}{t+b} \quad \text{for some parameters } a > 0, b \geq 0$$

### Extension: projected subgradient

When solving the dual problem, we need to enforce  $\lambda \succeq 0$

Can be done by projected subgradient method

$$x(t+1) = [x(t) - \alpha(t)h(x(t))]^+$$

where  $[\cdot]^+$  denotes projection on the positive orthant

Same steplength rules and convergence results as simple method

### Example: piecewise linear minimization

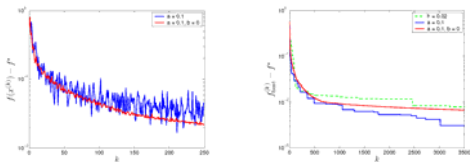
Consider the problem

$$\text{minimize } f(x) = \max_{k=1, \dots, K} \{a_k^T x + b_k\}$$

**Quiz:** determine the associated subgradient update

### Example: PWL minimization cont'd

Convergence of subgradient method for different stepsize rules

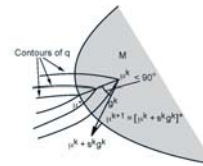


**Practical observation:** convergence is slow (in the limit)

### Subgradient method

Not a descent method

- function value iterates may increase
- distance to optimal set guaranteed to decrease (if stepsize small enough)



### Subgradients and supgradients

Similarly, a supgradient of a function  $g$  at  $x$  is any vector  $h$  such that

$$g(y) \leq g(x) + h^T(x)(y - x) \quad \text{for all } y$$

(supgradients give global overestimators)

For concave maximization, subgradient method reads

$$x^{(k+1)} = [x^{(k)} + \alpha^{(k)}h^{(k)}]^+$$

Many texts do not make a difference between sup and subgradients.

### Example: supgradient of dual function

Let  $x_\lambda$  be a minimizer for the Lagrangian, i.e.,

$$x_\lambda = \arg \min_{x \in X} L(x, \lambda) = \arg \min_{x \in X} \left\{ f_0(x) + \sum_{i=1}^m \lambda_i f_i(x) \right\}$$

then  $(f_1(x_\lambda) \dots f_m(x_\lambda))$  is a supgradient of the dual function at  $\lambda$

$$g(\bar{\lambda}) \leq g(\lambda) + \sum_{i=1}^m (\bar{\lambda}_i - \lambda_i) f_i(x_\lambda) \quad \text{for all } \bar{\lambda} \in \mathbb{R}^m$$

**Proof:**

**Observation:** subgradient for free when evaluating dual function!

## When is the dual differentiable?

**Fact:** If  $f_0$  is strictly convex,  $f_1, \dots, f_m$  linear, then  $g$  is differentiable (much weaker formulations exist)

## Disposition

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## Basic idea

Optimization problem in many variables

$$\begin{aligned} &\text{minimize} && \sum_{k=1}^n f_{0k}(x_k) \\ &\text{subject to} && \sum_{k=1}^n f_{ik}(x_k) \leq 0 \quad i = 1, \dots, m \\ &&& x_k \in X_k \quad i = 1, \dots, n \end{aligned}$$

with few coupling constraints (i.e.,  $n$  is large,  $m$  is small)

Introducing dual variables for coupling constraints only

$$\begin{aligned} g(\lambda) &= \inf_x \sum_k f_{0k}(x_k) + \sum_i \lambda_i \sum_k f_{ik}(x_k) \\ &= \sum_{k=1}^n \inf_{x \in X_k} f_{0k}(x_k) + \sum_i \lambda_i f_{ki}(x_k) \end{aligned}$$

Dual function separable – easy to evaluate if  $X_k$  has simple structure

## Solving the dual

Dual problem can be solved with projected supgradient method

$$\lambda_i(t+1) = \left[ \lambda_i(t) + \alpha(t) \sum_k f_{ik}(x_k^*(t)) \right]^+$$

At each iteration,  $g(\lambda)$  underestimates optimal objective

Convergence in the limit

## Example: resource allocation

$$\begin{aligned} &\text{minimize} && -\sum_{k=1}^n \log(x_k + \alpha_k) \\ &\text{subject to} && x_k \geq 0, \quad \mathbf{1}^T x \leq 1 \end{aligned}$$

Introduce Lagrange multiplier  $\lambda$  for coupling constraint  $\mathbf{1}^T x \leq 1$

Dual function

$$g(\lambda) = \sum_k \left( \inf_{x_k \geq 0} -\log(x_k + \alpha_k) + \lambda x_k \right) - \lambda$$

Inner problems solved by

$$x_k^* = \max\{0, 1/\nu - \alpha_k\}$$

Subgradient method:  $\lambda(t+1) = [\lambda(t) + \alpha(t)(\sum x_k - 1)]^+$

- Increase multipliers when resource overutilized, decrease if underutilized

## Pricing or tax interpretation

Operation of enterprise subject to resource constraints

$$\begin{aligned} &\text{minimize} && \sum_i f_i(x_i) \\ &\text{subject to} && \sum_i x_i \leq x_{\text{tot}} \end{aligned}$$

$f_i(x_i)$  cost of operating system  $i$  using  $x_i$  amount of resources

Dual function

$$g(\lambda) = \sum_i \inf_{x_i} f_i(x_i) + \lambda x_{\text{tot}}$$

Interpret  $\lambda$  as resource price: systems operate to minimize total cost

Subgradient method:

- increase resource prices when resource is overutilized,
- decrease resource prices when resource is underutilized

## Problems and challenges

Slow convergence of subgradient method

- cutting-plane methods faster, but computationally more intensive

Primal iterates  $x^*(t)$  not necessarily feasible

- coupling constraints not enforced in dual formulation
- need structure, or heuristic, to produce (suboptimal) primal solutions

## Disposition

- Lagrange duality review
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## Distributed optimization

In some cases, decomposition schemes reveal solutions that rely on

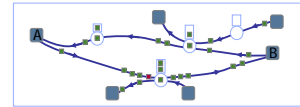
- distributed optimization (in local variables)
- and coordination mechanisms (dual variable updates/prices)

A powerful methodology for finding decentralized solutions!

We will simply exemplify this on a model of Internet congestion control.

## Internet congestion control

Users contend for limited network resources



Congestion control: helps share resources fairly and efficiently

- When router buffer fills up, packets are marked (or dropped)
- Users react to congestion signals by reducing transmission rate

Interplay between control mechanisms in routers and end nodes

## Internet Congestion Control

Current Internet

- source algorithm is some Transmission Control Protocol (TCP)
- link algorithm is some Active Queue Management (AQM) scheme

Two types of fundamental studies

- equilibrium properties (fairness, throughput, ...) – via cvx. opt.
- dynamical properties (stability, convergence) – control theory

Our focus: equilibrium properties using a dual decomposition approach

## Notation

Network with  $L$  links of finite capacities  $c_l$ ,  $l = 1, \dots, L$

Capacity shared by  $P$  source-destination pairs sending at rate  $s_p$

Sending rates must satisfy capacity constraints

$$\sum_p r_{lp} s_p \leq c_l \quad \text{where } r_{lp} = \begin{cases} 1 & \text{if source } s_p \text{ sends data across link } l \\ 0 & \text{otherwise} \end{cases}$$

Capacity constraints on vector form

$$r_l^T s \leq c_l \quad l = 1, \dots, L$$

## The utility maximization problem

Optimal network operation from solving utility maximization problem

$$\begin{aligned} & \text{maximize} && \sum_{p=1}^P u_p(s_p) \\ & \text{subject to} && r_l^T s \leq c_l \quad l = 1, \dots, L \\ & && s_p \geq 0 \quad p = 1, \dots, P \end{aligned}$$

where  $u_p$  is a strictly concave *utility function*

**Key observation:** the dual problem admits decentralized solution

- links update congestion measures (Lagrange multipliers)
- sources react to congestion signals, update source rates

Gives insight into current (and future) congestion control schemes!

## A dual approach

Lagrangian of utility maximization problem

$$L(x, \lambda) = \left\{ \sum_{p=1}^P u_p(s_p) - \sum_{l=1}^L \lambda_l (r_l^T s - c_l) \mid s_p \geq 0, p = 1, \dots, P \right\}$$

Dual function

$$\begin{aligned} g(\lambda) &= \max_{s \geq 0} \sum_{p=1}^P u_p(s_p) - \sum_{l=1}^L \lambda_l (r_l^T s - c_l) = \dots = \\ &= \sum_{p=1}^P \max_{s_p \geq 0} u_p(s_p) - \left( \sum_{l \in \mathcal{L}(p)} \lambda_l \right) s_p + \sum_{l=1}^L \lambda_l c_l = \sum_{p=1}^P B_p(q_p) + \sum_{l=1}^L \lambda_l c_l \end{aligned}$$

**Observation:** sources can deduce optimal rate from  $q_p = \sum_{l \in \mathcal{L}(p)} \lambda_l$

## Equilibrium rates and utilities

Sources maximize utilities minus "resource cost"

$$\max_{s_p \geq 0} u_p(s_p) - s_p \sum_{l \in \mathcal{L}(p)} \lambda_l = \max_{s_p \geq 0} u_p(s_p) - q_p s_p$$

Equilibrium rates from first-order conditions

$$u_p'(s_p^*) - q_p = 0 \Rightarrow s_p^* = \max \left\{ 0, (u_p')^{-1}(q_p) \right\}$$

Can also be used to deduce utility function from equilibrium  $(s_p^*, q_p^*)$

## Solving the dual problem

Apply (sub-)gradient method to compute optimal multipliers

$$\lambda_l^{(k+1)} = \left[ \lambda_l^{(k)} + \alpha^{(k)} \left( \sum_{p \in \mathcal{P}(l)} s_p - c_l \right) \right]^+$$

**Observation:** links update multipliers based on local excess capacity

Common to use constant step-length

## A synchronous gradient projection

**Link algorithm** (carried out at each link at each time  $k=1, \dots, K$ )

1. Compute the requested aggregate data rate  $\sum_{p \in \mathcal{P}(l)} s_p$
2. Update its price  $\lambda_l^{(k+1)} = [\lambda_l^{(k)} + \alpha^{(k)} (\sum_{p \in \mathcal{P}(l)} s_p - c_l)]^+$
3. Communicate new price to all sources  $p$  that use link  $l$

**Source algorithm** (executed at each source)

1. Receive sum of link prices along its path  $q_p = \sum_{l \in \mathcal{L}(p)} \lambda_l$
2. Choose new transmission rate  $s_p^{(t+1)} = s^*(q_p^{(t)})$
3. Communicate new rate to all links in path

## Primal and dual decomposition

Consider the problem

$$\begin{aligned} & \text{minimize} && f_x(x) + f_y(y) \\ & \text{subject to} && g_x(x) \leq r_x \quad g_y(y) \leq r_y \\ & && r_x + r_y \leq r_{\text{tot}} \end{aligned}$$

Dual decomposition: relax resource constraint, solve subproblems

$$\begin{aligned} & \text{minimize} && f_x(x) + \lambda r_x & \text{minimize} && f_y(y) + \lambda r_y \\ & \text{subject to} && g_x(x) \leq r_x & \text{subject to} && g_y(y) \leq r_y \end{aligned}$$

Coordinate by subgradient method (that finds optimal resource price)

## The primal approach

Re-write problem as

$$\text{minimize } p_x(r) + p_y(r)$$

where

$$p_x(r_x) = \inf \{f_x(x) \mid g_x(x) \leq r\}$$

$$p_y(r_x) = \inf \{f_y(y) \mid g_y(y) \leq r_{\text{tot}} - r\}$$

Recall:  $p_x(r_x)$  is convex, a subgradient is given by  $-\lambda_x$

Coordinator updates resource allocation (rather than prices)

- all iterates are feasible

## Primal decomposition

A perfect research paper presentation topic!

Reference: "Notes on Decomposition Methods", Boyd, Xiao and Mutapcic

## Summary

- subgradients and the subgradient method
- solving the primal via the dual
- dual decomposition for efficient optimization
- distributed optimization example: duality model of TCP