## SF2812 Applied linear optimization, final exam Friday January 152010 8.00-13.00

Examiner: Anders Forsgren, tel. 7907127.
Allowed tools: Pen/pencil, ruler and eraser.
Solution methods: Unless otherwise stated in the text, the problems should be solved by systematic methods, which do not become unrealistic for large problems. If you use methods other than what have been taught in the course, you must explain carefully.
Note! Personal number must be written on the title page. Write only one exercise per sheet. Number the pages and write your name on each page.
22 points are sufficient for a passing grade. For $20-21$ points, a completion to a passing grade may be made within three weeks from the date when the results of the exam are announced.

1. Consider the linear program $(L P)$ defined as

$$
\begin{array}{lll} 
& \min & x_{1}+2 x_{2} \\
(L P) & \text { då } & x_{1}+3 x_{2}=1 \\
& & x_{1} \geq 0, x_{2} \geq 0 .
\end{array}
$$

Consider a barrier transformation of $(L P)$ for a given positive barrier parameter $\mu$.
(a) For the given $\mu$, formulate the primal-dual system of nonlinear equations corresponding to the problem above. Use the fact that the problem is small to calculate an explicit expression for the solution $x(\mu), y(\mu)$ and $s(\mu)$ to these nonlinear equations.
(b) Show that the calculated $x(\mu), y(\mu)$ and $s(\mu)$ converge to an optimal solution of $(L P)$ and corresponding dual problem as $\mu \rightarrow 0$.
2. Consider a binary knapsack problem $(K P)$ defined as

$$
\begin{array}{ll}
\text { minimize } & -\sum_{j=1}^{n} c_{j} x_{j} \\
\text { subject to } & -\sum_{j=1}^{n} a_{j} x_{j} \geq-b, \\
& x_{j} \in\{0,1\}, \quad j=1, \ldots, n,
\end{array}
$$

where $a \geq 0, c \geq 0$ and $b \geq 0$.
(a) Give an explicit expression for the objective function $\varphi(\lambda)$ of the dual problem $(D)$ arising when the constraint $-\sum_{j=1}^{n} a_{j} x_{j} \geq-b$ is relaxed by Lagrangian relaxation.
(b) For a given $\lambda \in \mathbb{R}$, find an explicit expression for a subgradient to $\varphi$ at $\lambda$. (2p)
(c) Assume that $n=3, a=\left(\begin{array}{ll}2 & 3\end{array}\right)^{T}, b=6$ and $c=\left(\begin{array}{lll}4 & 5 & 6\end{array}\right)^{T}$. Illustrate the dual problem graphically. From the figure calculate the optimal solution and the optimal objective value of the dual problem. Solve the small (KP) by inspection, and determine the duality gap. ..................................... (5p)
3. Consider the stochastic program $(P)$ given by
(P)

$$
\begin{array}{cl}
\operatorname{minimize} & c^{T} x \\
\text { subject to } & A x=b  \tag{P}\\
& T(\omega) x=h(\omega) \\
& x \geq 0
\end{array}
$$

where $\omega$ is a stochastic variable and $T(\omega) x=h(\omega)$ is to be interpreted as an "informal" stochastic constraint. Assume that $\omega$ takes on a finite number of values $\omega_{1}, \ldots, \omega_{N}$ with corresponding probabilities $p_{1}, \ldots, p_{N}$. Let $T_{i}$ denote $T\left(\omega_{i}\right)$ and let $h_{i}$ denote $h\left(\omega_{i}\right)$.
(a) Explain how the deterministically equivalent problem

$$
\begin{array}{ll}
\text { minimize } & c^{T} x+\sum_{i=1}^{N} p_{i} q_{i}^{T} y_{i} \\
\text { subject to } & A x=b, \\
& T_{i} x+W y_{i}=h_{i}, \quad i=1, \ldots, N \\
& x \geq 0, \\
& y_{i} \geq 0, \quad i=1, \ldots, N
\end{array}
$$

arises. (We assume, for simplicity, "fix compensation", i.e., $W$ does not depend on $i$.)
(b) Define $V S S$ in terms of suitable optimization problems. .................... (2p)
(c) Define EVPI in terms of suitable optimization problems.
4. Consider a transportation problem (PTP) defined as

$$
\begin{array}{ll}
\operatorname{minimize} & \sum_{i=1}^{4} \sum_{j=1}^{3} c_{i j} x_{i j} \\
\text { subject to } & \sum_{j=1}^{3} x_{i j}=a_{i}, \quad i=1,2,3,4 \\
& \sum_{i=1}^{4} x_{i j}=b_{j}, \quad j=1,2,3 \\
& x_{i j} \geq 0, \quad i=1,2,3,4, \quad j=1,2,3
\end{array}
$$

with corresponding dual problem $(D T P)$ given by

$$
\begin{array}{rll} 
& \text { maximize } & \sum_{i=1}^{4} a_{i} u_{i}+\sum_{j=1}^{3} b_{j} v_{j} \\
(D T P) & \text { subject to } \quad & u_{i}+v_{j}+s_{i j}=c_{i j}, \quad i=1,2,3,4, \quad j=1,2,3, \\
& s_{i j} \geq 0, \quad i=1,2,3,4, \quad j=1,2,3
\end{array}
$$

where

$$
C=\left(\begin{array}{rrr}
8 & 6 & 10 \\
9 & 1 & 10 \\
1 & 3 & 2 \\
9 & 5 & 10
\end{array}\right), \quad a=\left(\begin{array}{l}
1 \\
2 \\
3 \\
2
\end{array}\right), \quad b=\left(\begin{array}{l}
2 \\
2 \\
4
\end{array}\right)
$$

AF has implemented a primal-dual interior method for linear programming. When solving this transportation problem he obtains

$$
\widehat{X}=\left(\begin{array}{ccc}
1 & 0 & 0  \tag{3p}\\
0 & 2 & 0 \\
\frac{1}{2} & 0 & \frac{5}{2} \\
\frac{1}{2} & 0 & \frac{3}{2}
\end{array}\right), \quad \widehat{u}=\left(\begin{array}{c}
9 \\
\frac{15}{2} \\
2 \\
10
\end{array}\right), \quad \widehat{v}=\left(\begin{array}{c}
-1 \\
-\frac{13}{2} \\
0
\end{array}\right), \quad \widehat{S}=\left(\begin{array}{ccc}
0 & \frac{7}{2} & 1 \\
\frac{5}{2} & 0 & \frac{5}{2} \\
0 & \frac{15}{2} & 0 \\
0 & \frac{3}{2} & 0
\end{array}\right) .
$$

(a) Show that $\widehat{X}$ is an optimal solution to (PTP).
(b) AF is a bit confused. He expected the optimal $X$ to have integer values only, as he is solving a transportation problem with $a$ and $b$ integer valued. Explain why it is not surprising that AF's interior solver has not produced an optimal $X$ with integer values only.
(c) Find, using $\widehat{X}$, two integer optimal solutions to (PTP).

Hint: The matrix $P$ given by

$$
P=\left(\begin{array}{rrr}
0 & 0 & 0 \\
0 & 0 & 0 \\
1 & 0 & -1 \\
-1 & 0 & 1
\end{array}\right)
$$

might be helpful.
(d) If (PTP) was solved by the simplex method, which of the three optimal solutions mentioned above might be obtained? Motivate your answer.
5. Consider a cutting-stock problem with the following data:

$$
W=11, \quad m=3, \quad w_{1}=3, \quad w_{2}=5, \quad w_{3}=9, \quad b=\left(\begin{array}{ccc}
60 & 50 & 40
\end{array}\right)^{T}
$$

Notation and problem statement are in accordance to the textbook. Given are rolls of width $W$. Rolls of $m$ different widths are demanded. Roll $i$ has width $w_{i}$, $i=1, \ldots, m$. The demand for roll $i$ is given by $b_{i}, i=1, \ldots, m$. The aim is to cut the $W$-rolls so that a minimum number of $W$-rolls are used.
(a) Solve the the LP-relaxed problem associated with the above problem. Start with the basic feasible solution associated with the three "pure" cut patterns $\left(\begin{array}{lll}3 & 0 & 0\end{array}\right)^{T},\left(\begin{array}{lll}0 & 2 & 0\end{array}\right)^{T}$ and $\left(\begin{array}{lll}0 & 0 & 1\end{array}\right)^{T}$. The subproblems that arise may be solved in any way, that need not be systematic.
(b) Determine a "near-optimal" solution to the original problem. Give a bound on the maximum deviation from the optimal value of the original problem... (2p)

