

## SF2812 Applied linear optimization, final exam Wednesday June 10 2015 14.00–19.00

Examiner: Anders Forsgren, tel. 08-790 71 27.

Allowed tools: Pen/pencil, ruler and eraser.

*Note!* Calculator is not allowed.

*Solution methods:* Unless otherwise stated in the text, the problems should be solved by systematic methods, which do not become unrealistic for large problems. Motivate your conclusions carefully. If you use methods other than what have been taught in the course, you must explain carefully.

*Note!* Personal number must be written on the title page. Write only one exercise per sheet. Number the pages and write your name on each page.

22 points are sufficient for a passing grade. For 20-21 points, a completion to a passing grade may be made within three weeks from the date when the results of the exam are announced.

**1.** Let (LP) be defined as

(LP) minimize 
$$c^T x$$
  
subject to  $Ax = b$ ,  
 $x \ge 0$ ,

where

$$A = \begin{pmatrix} 2 & 3 & 1 & 0 & 0 \\ -1 & 2 & 0 & 1 & 0 \\ 2 & 1 & 0 & 0 & 1 \end{pmatrix}, \quad b = \begin{pmatrix} 4 \\ 1 \\ 5 \end{pmatrix} \quad \text{and} \quad c = \begin{pmatrix} -8 & 2 & 0 & 2 & -2 \end{pmatrix}^T.$$

(a) A person named AF has used GAMS to model and solve this problem. The GAMS input file can be found at the end of the exam. Unfortunately, AF has lost the GAMS output file. He does have a partial GAMS output file, which reads:

EQU cons	constrain	ts						
LOWER	LEVEL	UPPER	MARGIN	AL				
	4.000							
i2 1.000 i3 5.000								
	L	OWER	LEVEL	UPPER	MARGINAL			
EQU objfı	ın	•	•	•	-1.000			
objfun objective function								
	L	OWER	LEVEL	UPPER	MARGINAL			
VAR objva	al ·	-INF -	12.000	+INF				
objval objective function value								
VAR x decision variables								

	LOWER	LEVEL	UPPER	MARGINAL
j1		2.000	+INF	
j2	•	•	+INF	3.000
jЗ	•	•	+INF	1.000
j4	•	3.000	+INF	
j5	•	1.000	+INF	

- **2.** Consider a binary knapsack problem (KP) defined as

(KP) minimize 
$$-\sum_{j=1}^{n} c_j x_j$$
  
subject to  $-\sum_{j=1}^{n} a_j x_j \ge -b,$   
 $x_j \in \{0,1\}, \quad j = 1, \dots, n,$ 

where  $a \ge 0, c \ge 0$  and  $b \ge 0$ .

- (b) For a given  $\lambda \in \mathbb{R}$ , find an explicit expression for a subgradient to  $\varphi$  at  $\lambda$ . (2p)
- **3.** Consider the linear program

(LP) minimize 
$$c^T x$$
  
subject to  $Ax = b$ ,  
 $x \ge 0$ ,

where

$$A = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & 2 & 3 & 4 \end{pmatrix}, \quad b = \begin{pmatrix} 10 \\ 20 \end{pmatrix}, \quad c = \begin{pmatrix} 2 & 3 & 3 & 3 \end{pmatrix}^{T}.$$

Assume that we want to solve (LP) using a primal-dual interior-point method. Assume further that we initially choose  $x^{(0)} = (4\ 3\ 2\ 1)^T$ ,  $y^{(0)} = (0\ 0)^T$ ,  $s^{(0)} = (1\ 2\ 3\ 4)^T$ , and  $\mu = 1$ . Here, y and s denote the dual variables.

- (b) The solution to the above system of linear equations is given by

$$\Delta x \approx \begin{pmatrix} 0.4588\\ -0.7000\\ 0.0235\\ 0.2176 \end{pmatrix}, \quad \Delta y \approx \begin{pmatrix} 1.5294\\ 0.3353 \end{pmatrix}, \quad \Delta s \approx \begin{pmatrix} -0.8647\\ -1.2000\\ -2.5353\\ -3.8706 \end{pmatrix}$$

4. Consider the stochastic program (P) given by

$$\begin{array}{ll} \mbox{minimize} & c^T x \\ (P) & \mbox{subject to} & Ax = b, \\ & & T(\omega)x = h(\omega), \\ & & x \geq 0, \end{array}$$

where  $\omega$  is a stochastic variable and  $T(\omega)x = h(\omega)$  is to be interpreted as an "informal" stochastic constraint. Assume that  $\omega$  takes on a finite number of values  $\omega_1, \ldots, \omega_N$  with corresponding probabilities  $p_1, \ldots, p_N$ . Let  $T_i$  denote  $T(\omega_i)$  and let  $h_i$  denote  $h(\omega_i)$ .

(a) Explain how the deterministically equivalent problem

minimize 
$$c^T x + \sum_{i=1}^N p_i q_i^T y_i$$
  
subject to  $Ax = b$ ,  
 $T_i x + W y_i = h_i, \quad i = 1, \dots, N,$   
 $x \ge 0,$   
 $y_i \ge 0, \quad i = 1, \dots, N,$ 

5. Consider a cutting-stock problem with the following data:

$$W = 11, \quad m = 3, \quad w_1 = 3, \quad w_2 = 5, \quad w_3 = 9, \quad b = \begin{pmatrix} 60 & 50 & 40 \end{pmatrix}^T.$$

Notation and problem statement are in accordance to the textbook. Given are rolls of width W. Rolls of m different widths are demanded, where roll i has width  $w_i$ ,  $i = 1, \ldots, m$ . The demand for roll i is given by  $b_i$ ,  $i = 1, \ldots, m$ . The aim is to cut the W-rolls so that a minimum number of W-rolls are used.

- (b) Determine a "near-optimal" solution to the original problem. Give a bound on the maximum deviation from the optimal value of the original problem...(2p)

Good luck!

GAMS file for Question 1:

```
sets
i rows
            / i1*i3 /
            / j1*j5 /;
j columns
table A(i,j) constraint matrix
       j1 j2 j3
                     j4
                          j5
        2 3
i1
                1
           2
       -1
                      1
i2
        2
             1
i3
                           1;
parameter b(i)
               4
         / i1
           i2
               1
           i3 5/;
parameter c(j)
         / j1 -8
           j2 2
           j4 2
           j5 -2 /;
variables
       objval objective function value
       x(j)
               decision variables;
positive variable x;
equations
       cons(i)
                      constraints
       objfun
                      objective function;
cons(i) .. sum(j,A(i,j)*x(j)) =e= b(i);
objfun .. sum(j,c(j)*x(j)) =e= objval;
model lpex / all /;
solve lpex using lp minimizing objval;
```