

Mathematical Foundation for Compressed Sensing

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Lecture 5, March 5, 2012

An outline for today

- A short summary from last lecture:

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- RIP of Random matrices (continued).

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Theorem 3.3: If \mathbf{A} fulfils the RIP property with constant $\delta_{2s} < \sqrt{2} - 1 \approx 0.412\dots$, then $\mathbf{A} \in \text{MP}_1(s)$.

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Best known: enough with $\frac{\delta_{2s}}{\sqrt{1-\delta_{2s}^2}} < \frac{1}{\sqrt{2}}$. Thus it is enough with $\delta_{2s} < \frac{1}{\sqrt{3}} \approx 0.577\dots$

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Lemma 3.5: If \mathbf{A} satisfies RIP properties and $\mathbf{Ax} = 0$, then

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Remark: With some effort the constant $\frac{\delta_{2s}}{1 - \delta_s}$ in Lemma 3.5 can be replaced by $\frac{\delta_{2s}}{\sqrt{1 - \delta_{2s}^2}}$.

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More precisely we get

$$m \geq \frac{s \log(2Ne/s) - \log 4}{2(\log(\sqrt{1 - \delta_s} + 4\sqrt{1 + \delta_s}) - \log(1 - \delta_s))}.$$

RIP for Gaussian Matrices

Let random matrix $\mathbf{A} = (A_{ij})_{1 \leq i \leq m, 1 \leq j \leq N}$ where each A_{ij} is an independent Gaussian random variable, i.e each $A_{ij} \in N(0, 1)$ with distribution function $\phi(t)$ where

$$\phi(t) = \frac{1}{\sqrt{2\pi}} e^{-t^2/2}.$$

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$$m > C \frac{s}{\delta^2} \log(N/(s\epsilon)) + \text{lower order terms},$$

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More precisely

$$m > \frac{2s \log(N/s) + 4s \log \log(N/s)}{\delta_s - \log(1 + \delta_s)} + \text{lower order terms}.$$

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Beginning of the proof of Theorem 4.5:

Let \mathbf{x} be a vector in R^N with length $\|\mathbf{x}\|_2 = 1$ and define the random variable $T = \|\mathbf{Ax}\|_2^2/m$.

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- Write a vector-distribution function of (Z_1, \dots, z_m)

$$\phi((t_1, \dots, t_m)) = (2\pi)^{-m/2} e^{-\sum_i t_i^2/2}.$$

- Set $\tau = \sum_i t_i^2/m$. It follows that T has distribution function

$$\Psi(\tau) = c_m \tau^{\frac{m}{2}-1} e^{-m\tau/2},$$

for some normalization constant c_m .

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- second derivative of $\Psi(\tau)$

$$\frac{d^2}{dx^2} \log \Psi(\tau) = -\left(\frac{m}{2} - 1\right) \frac{1}{\tau^2}.$$

We conclude

Lemma 5.1 $\Psi(\tau)$ has maximum value

$$\Psi_{max} = \Psi_{max,m} \sim \sqrt{m}/\sqrt{2\pi} \text{ at } \tau_m = 1 - \frac{2}{m}.$$

Furthermore

$$\int_{|\tau-1|>\delta} \Psi(\tau) d\tau \sim \frac{\sqrt{2}e^{-\frac{m}{2}(\delta-\log(1+\delta))}}{\sqrt{m\pi}\delta}.$$

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And we have

$$\delta - \log(1 + \delta) < -\delta - \log(1 - \delta).$$

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This implies the integral over $0 < \tau < 1 - \delta$ neglectible for large m .

Thus for one fixed point probability for the RIP estimate does **not** hold, is of magnitude

$$\frac{\sqrt{2}e^{-\frac{m}{2}(\delta-\log(1+\delta))}}{\sqrt{m\pi\delta}},$$

for large m .

We summarise in the

Lemma 5.1 Let \mathbf{x} any fixed vector in \mathbb{R}^N and let \mathbf{A} be a Gaussian random matrix then we have the RiP estimate

$$| \|\mathbf{Ax}\|_2^2 - \|\mathbf{x}\|_2^2 | \leq \delta \|\mathbf{x}\|_2^2$$

with probability $1 - \epsilon$ where

$$\epsilon = \frac{\sqrt{2}e^{-\frac{m}{2}(\delta - \log(1+\delta))}}{\sqrt{m\pi}\delta}$$

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Corollary: This gives the RIP estimate with constant $\delta_s = s\delta$ with probability larger than $1 - \epsilon$ provided

$$m > \frac{2}{\delta/s - \log(1 - \delta/s)} (\log(N^2/\epsilon) + c) \sim \frac{4s^2}{\delta^2} (\log(N^2/\epsilon) + c).$$

We have used that $X(s)$ is contained in the ball $B(\sqrt{s}) = \{\mathbf{x} : \|\mathbf{x}\|_1 \leq \sqrt{s}\}$ which is the convex hull of the points $\mathbf{x}_{\pm i} = \pm\sqrt{s}\mathbf{e}_i$.

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When taking the convex hull we lost a factor \sqrt{s}
void this by using more points.

Recall the Definitions

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Recall the Definitions

Let $0 < r < 1$, and let $N_s(r)$ be the maximal number of points $\{\mathbf{x}_i\}$ in a set in the unit ball $B_s = \{\mathbf{x} \in \mathbb{R}^m : \|\mathbf{x}\|_2 \leq 1\}$ in \mathbb{R}^s such that

$$\|\mathbf{x}_i - \mathbf{x}_j\|_2 > r \text{ for all } i \neq j.$$

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Lemma 5.4: Let $\{\mathbf{x}_i\}_1^M$ a set with maximal number $M = N_s(r)$ of points in the unit ball B_s .

Then the closed convex hull W of $\{\mathbf{x}_i\}$ contains $(1 - r)B_s = \{x : \|\mathbf{x}\| \leq 1 - r\}$.

Proof: Let \mathbf{z} be a point in the unit ball B_s that is not in W .

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stated otherwise there: is a linear functional $l(x)$ on R^N such that $l(\mathbf{z}) > 1$ but $l(x) \leq 1$ for all $x \in W$.

Proof of Lemma 5.4 continued

Let \mathbf{z}_0 be point on the hyperplan with minimal distance to the origin and let $\mathbf{x}_{M+1} = \mathbf{z}_0 / \|\mathbf{z}\|_2$.

Proof of Lemma 5.4 continued

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Then for any $\mathbf{x}_i, 1 \leq i \leq M$

$$\|\mathbf{x}_i - \mathbf{x}_{M+1}\|_2 \geq \|\mathbf{z}_0 - \mathbf{x}_{M+1}\|_2 = 1 - \|\mathbf{z}_0\|_2 \geq 1 - \|\mathbf{z}\|_2.$$

Proof of Lemma 5.4 continued

Let \mathbf{z}_0 be point on the hyperplan with minimal distance to the origin and let $\mathbf{x}_{M+1} = \mathbf{z}_0 / \|\mathbf{z}\|_2$.

Then for any $\mathbf{x}_i, 1 \leq i \leq M$

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Finishing the proof of Theorem 4.5:

Recall Theorem 4.5: There is a constant $C > 0$ such for any $\epsilon > 0, \delta > 0$ if we let

$$m > C \frac{s}{\delta^2} \log(Ne/(s\epsilon)).$$

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i.e points \mathbf{x} supported on S an with norm $\|\mathbf{x}\|_2 = 1 - r$.

We will use the RIP estimate for all the points: $\mathbf{x}_{S,i}, 1 \leq i \leq M$, $\mathbf{x}_{S,i} + \mathbf{x}_{S,j}, i \neq j$ and $\mathbf{x}_{S,i} - \mathbf{x}_{S,j}, i \neq j$ according to Lemma 5.1.

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Theorem 4.5 There is a constant $C > 0$ such that for any $\epsilon > 0$, and $0 < \delta \leq 1$, the following holds: If

$$m > \frac{2s \log(N/(s\epsilon)) + 4s \log \log(N/s)}{\delta_s - \log(1 + \delta_s)} + \text{lower order terms}$$

then the Gaussian $m \times N$ matrix \mathbf{A} (as above) satisfy RIP constant $\delta_s = \delta$ with a probability not less than $1 - \epsilon$.

Corollary There exist $m \times N$ matrices RIP constant δ_s if

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Construction of matrices satisfyin RIP

Construction of RIP matrices We want do a construction of an $m \times N$ matrix \mathbf{A} statistfying the RIP property

Theorem 5.5 There is a constant $C > 0$ such that if $\delta_s > 0$ and

$$m > C \frac{s^2}{\delta_s^2} \log N.$$

then it is possible to construct an $m \times N$ matrix \mathbf{A} statistfying the RIP property with constant δ_s .

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Proof is done by construction:

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- 6 Choose the new column A_n on the unite sphere, but outside each such plate.

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Proof of RIP property of constructed matrix **A**:

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- let $W = \text{span}\{A_{n_j}, j < k\}$ then $\mathbf{u} = A_{n_k}$ is uniquely decomposed

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- set

$$\mathbf{n}_{j_1, \dots, j_k} = \frac{\mathbf{u}_{W^\perp}}{\|\mathbf{u}_{W^\perp}\|_2} \text{ and } \mathbf{v}_{j_1, \dots, j_k} = \mathbf{u} - \mathbf{n}_{j_1, \dots, j_k}$$



$$\mathbf{U} = (\mathbf{n}_{j_1, \dots, j_k})_{k=1}^s$$

will be an orthonormal matrix with orthonormal columns and



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- if \mathbf{x} is supported on $\{j_1, \dots, j_s\}$
will we get $\mathbf{Ax} = \mathbf{Ux} + \mathbf{Vx}$ and

$$\|\mathbf{Ux}\|_2 - \|\mathbf{Vx}\|_2 \leq \|\mathbf{Ax}\|_2 \leq \|\mathbf{Ux}\|_2 + \|\mathbf{Vx}\|_2$$

- Since $\|\mathbf{U}\mathbf{x}\|_2 = \|\mathbf{x}\|_2$ will we get

$$|\|\mathbf{Ax}\|_2^2 - \|\mathbf{x}\|_2^2| \leq 2\|\mathbf{x}\|_2\|\mathbf{Vx}\|_2 + \|\mathbf{Vx}\|_2^2.$$

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By the triangle inequality we get

$$\|\mathbf{Vx}\|_2 \leq \sqrt{2(1-r)}\|\mathbf{x}\|_1 \leq \sqrt{2s(1-r)}\|\mathbf{x}\|_2.$$

- Since $\|\mathbf{U}\mathbf{x}\|_2 = \|\mathbf{x}\|_2$ will we get

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By the triangle inequality we get

$$\|\mathbf{V}\mathbf{x}\|_2 \leq \sqrt{2(1-r)}\|\mathbf{x}\|_1 \leq \sqrt{2s(1-r)}\|\mathbf{x}\|_2.$$

Thus

$$|\|\mathbf{A}\mathbf{x}\|_2^2 - \|\mathbf{x}\|_2^2| \leq (2\sqrt{2s(1-r)} + 2s(1-r))\|\mathbf{x}\|_2^2.$$

- Finally we choose r such that $2\sqrt{2s(1-r)} + 2s(1-r) \leq \delta_s$

Careful estimate of the covering and the conditions for m in Lemma 5.6 will give us the estimate how large m needs to be in Theorem 5.5