

**KTH Mathematics** 

## SF2822 Applied nonlinear optimization, final exam Saturday December 20 2008 8.00–13.00

Examiner: Anders Forsgren, tel. 790 71 27.

Allowed tools: Pen/pencil, ruler and eraser; plus a calculator provided by the department. Solution methods: Unless otherwise stated in the text, the problems should be solved by systematic methods, which do not become unrealistic for large problems. If you use methods other than what have been taught in the course, you must explain carefully.

*Note!* Personal number must be written on the title page. Write only one exercise per sheet. Number the pages and write your name on each page.

22 points are sufficient for a passing grade. For 20-21 points, a completion to a passing grade may be made within three weeks from the date when the results of the exam are announced.

1. Consider the nonlinear optimization problem (NLP) defined as

(*NLP*) minimize 
$$e^{x_1} + \frac{1}{2}(x_1 + x_2 - 4)^2 + (x_1 - x_2)^2$$
  
subject to  $-(x_1 - 3)^2 - x_2^2 + 9 \ge 0.$ 

You have obtained a printout from a sequential quadratic programming solver for this problem. The initial point is  $x = (0 \ 0)^T$  and  $\lambda = 0$ . Six iterations, without linesearch, have been performed. The printout, where the floating point numbers are given with four decimal places, reads:

$\operatorname{It}$	$x_1$	$x_2$	$\lambda$	$\ \nabla f(x) - \nabla g(x)\lambda\ $
0	0	0	0	5.0000
1	0	1.3333	-0.7222	1.9259
2	1.0117	2.9430	-0.5596	1.1891
3	0.7084	2.1240	-0.4384	0.2028
4	0.7990	2.0422	-0.3294	0.0335
5	0.7983	2.0378	-0.3227	0.0001
6	0.7984	2.0378	-0.3227	0.0000

- (c) The printout corresponds to a problem related to (NLP). Which one? ... (2p)

Note: According to the convention of the book we define the Lagrangian  $\mathcal{L}(x,\lambda)$  as  $\mathcal{L}(x,\lambda) = f(x) - \lambda^T g(x)$ , where f(x) the objective function and g(x) is the constraint function, where the inequality constraint is written as  $g(x) \ge 0$ .

**2.** Consider the quadratic program (QP) defined by

(QP) minimize 
$$3x_1^2 - 2x_1x_2 + 3x_2^2 - 24x_1 - 8x_2$$
  
subject to  $-x_1 - x_2 \ge -5$ ,  
 $x_1 \ge 0$ ,  
 $x_2 \ge 0$ .

The problem may be illustrated geometrically in the figure below,



- (b) Assume that the constraint  $x_2 x_1 \ge -3$  is added to (QP), so that we obtain the problem (QP') according to

$$(QP') \qquad \begin{array}{ll} \text{minimize} & 3x_1^2 - 2x_1x_2 + 3x_2^2 - 24x_1 - 8x_2 \\ \text{subject to} & -x_1 - x_2 \ge -5, \\ & x_1 \ge 0, \\ & x_2 \ge 0, \\ & x_2 - x_1 \ge -3. \end{array}$$

3. Consider the nonlinear programming problem

(P) 
$$\begin{array}{c} \text{minimize} \quad f(x) \\ \text{subject to} \quad g(x) \ge 0, \end{array}$$

where  $f : \mathbb{R}^n \to \mathbb{R}$  and  $g : \mathbb{R}^n \to \mathbb{R}^m$  are continuously differentiable.

A barrier transformation of (P) for a fixed positive barrier parameter  $\mu$  gives the problem

$$(P_{\mu})$$
 minimize  $f(x) - \mu \sum_{i=1}^{m} \ln(g_i(x)).$ 

Show that the first-order necessary optimality conditions for  $(P_{\mu})$  are equivalent to the system of nonlinear equations

$$abla f(x) - 
abla g(x)\lambda = 0,$$
  
 $g_i(x)\lambda_i - \mu = 0, \quad i = 1, \dots, m,$ 

## 4. Consider a nonlinear programming problem (NLP) defined by

(*NLP*) minimize 
$$e^{x_1} + \frac{1}{2}x_1^2 + x_1x_2 + \frac{1}{2}x_2^2 + x_3^2 - 2x_1 - x_2 - 3x_3$$
  
(*NLP*) subject to  $x_1^2 + x_2^2 + x_3^2 + x_1 + x_2 - x_3 = 0,$   
 $-x_1^2 - x_2^2 - x_3^2 \ge -2.$ 

Let  $\tilde{x} = (0 \ 0 \ 1)^T$ .

- (a) Show that  $\tilde{x}$  fulfils the first-order necessary optimality conditions for (NLP). (3p)

- 5. For a given symmetric  $n \times n$ -matrix W, the so-called *two-way partitioning problem* may be posed as

$$\begin{array}{ll} \underset{x \in \mathbb{R}^n}{\text{minimize}} & x^T W x\\ \text{subject to} & x_{ii}^2 = 1, \quad i = 1, \dots, n. \end{array}$$

which is a nonconvex problem in general. By Lagrangian relaxation, one may obtain the dual problem as a semidefinite program  $(SDP_1)$  on the form

$$(SDP_1) \qquad \begin{array}{ll} \underset{y \in \mathbb{R}^n}{\text{minimize}} & -e^T y \\ \text{subject to} & \text{diag}(y) \preceq W, \end{array}$$

Alternatively, one may reformulate the original problem as

$$\begin{array}{ll} \underset{X \in \mathcal{S}^n}{\minininize} & \operatorname{trace}(WX)\\ \text{subject to} & X_{ii} = 1, \quad i = 1, \dots, n,\\ & \operatorname{rank}(X) = 1, \quad X \succeq 0, \end{array}$$

where  $S^n$  denotes the set of symmetric  $n \times n$  matrices. If the rank constraint is relaxed, one obtains the semidefinite program  $(SDP_2)$  given by

$$(SDP_2) \qquad \begin{array}{ll} \underset{X \in \mathcal{S}^n}{\text{minimize}} & \text{trace}(WX) \\ \text{subject to} & X_{ii} = 1, \quad i = 1, \dots, n, \\ & X \succeq 0. \end{array}$$

What is the relation between  $(SDP_1)$  and  $(SDP_2)$ ? .....(10p) Note: You need not consider the original problem, but you may rather consider  $(SDP_1)$  and  $(SDP_2)$  only.

*Hint 1:* Relate  $(SDP_1)$  and  $(SDP_2)$  via duality. Recall that a primal-dual pair of semidefinite programs are given by

$$(PSDP) \qquad \begin{array}{ll} \underset{X \in \mathcal{S}^n}{\text{minimize}} & \text{trace}(CX) \\ \text{subject to} & \text{trace}(A_iX) = b_i, \quad i = 1, \dots, m, \\ & X \succeq 0, \end{array}$$

and

$$(DSDP) \qquad \begin{array}{l} \underset{y \in \mathbb{R}^m}{\text{maximize}} \quad \sum_{i=1}^m b_i y_i \\ \text{subject to} \quad \sum_{i=1}^m A_i y_i \leq C. \end{array}$$

*Hint 2:* Note that if M is a symmetric  $n \times n$ -matrix and u is an n-vector, then  $u^T M u = \text{trace}(M u u^T)$ .

Good luck!



Figure for Exercise 2a:

Figure for Exercise 2b:

