

## SF2812 Applied linear optimization, final exam Wednesday January 16 2008 8.00–13.00

Examiner: Anders Forsgren, tel. 790 71 27.

Allowed tools: Pen/pencil, ruler and rubber; plus a calculator provided by the department.

Solution methods: Unless otherwise stated in the text, the problems should be solved by systematic methods, which do not become unrealistic for large problems. If you use methods other than what have been taught in the course, you must explain carefully.

*Note!* Personal number must be written on the title page. Write only one exercise per sheet. Number the pages and write your name on each page.

22 points are sufficient for a passing grade. For 20-21 points, a completion to a passing grade may be made within three weeks from the date when the results of the exam are announced.

1. Let (LP) be defined as

$$(LP) \qquad \begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax = b, \\ & x \geq 0, \end{array}$$

where

$$A = \begin{pmatrix} 2 & 4 & -1 & 0 & 0 \\ 1 & 1 & 0 & -1 & 0 \\ 3 & 1 & 0 & 0 & -1 \end{pmatrix}, \quad b = \begin{pmatrix} 2 \\ 1 \\ 1 \end{pmatrix} \quad \text{and} \quad c = \begin{pmatrix} 8 & 7 & -1 & -1 & 0 \end{pmatrix}^T.$$

Assume that we want to solve (LP) by a primal-dual interior method.

Let the initial point be given by

$$x = \begin{pmatrix} 1 \\ 2 \\ 8 \\ 2 \\ 4 \end{pmatrix}, \quad y = \begin{pmatrix} 1.1 \\ 1.1 \\ 1 \end{pmatrix}, \quad s = \begin{pmatrix} 1.7 \\ 0.5 \\ 0.1 \\ 0.1 \\ 1 \end{pmatrix}.$$

- **2.** Let (LP) be defined as

(LP) minimize 
$$c^T x$$
  
subject to  $Ax = b$ ,  $x > 0$ ,

where

$$A = \begin{pmatrix} 2 & 4 & -1 & 0 & 0 \\ 1 & 1 & 0 & -1 & 0 \\ 3 & 1 & 0 & 0 & -1 \end{pmatrix}, \quad b = \begin{pmatrix} 2 \\ 1 \\ 1 \end{pmatrix} \quad \text{and} \quad c = \begin{pmatrix} 8 & 7 & -1 & -1 & 0 \end{pmatrix}^T.$$

A person named AF claims that he has obtained  $x^* = (0\ 1\ 2\ 0\ 0)^T$  as an optimal solution to (LP) by the simplex method. However, he is a bit confused, since he would expect an optimal solution to have three positive components. Your task is to clarify AF's confusion.

- 3. Consider the binary integer programming problem (IP) given by

(IP) minimize 
$$-5x_1 - 7x_2 - 10x_3$$
  
subject to  $-3x_1 - 6x_2 - 7x_3 \ge -8$ ,  
 $-x_1 - 2x_2 - 3x_3 \ge -3$ ,  
 $x_j \in \{0, 1\}, \quad j = 1, \dots, n$ .

Assume that the constraint  $-3x_1 - 6x_2 - 7x_3 \ge -8$  is relaxed by Lagrangian relaxation for a nonnegative multiplier u.

- **4.** Consider the integer program (IP) defined by

$$(IP) \qquad \begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \geq b, \\ & Cx \geq d, \\ & x \geq 0, \quad x \text{ integral.} \end{array}$$

Let  $z_{IP}$  denote the optimal value of (IP).

Associated with (IP) we may define the dual problem (D) as

(D) 
$$\begin{array}{c} \text{maximize} \quad \varphi(u) \\ \text{subject to} \quad u \ge 0, \end{array}$$

where  $\varphi(u) = \min\{c^T x + u^T (b - Ax) : Cx \ge d, x \ge 0 \text{ integral}\}$ . Let  $z_D$  denote the optimal value of (D).

Let (LP) denote the linear program obtained from (IP) by relaxing the integrality requirement, i.e.,

$$(LP) \qquad \begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \geq b, \\ & Cx \geq d, \\ & x \geq 0. \end{array}$$

Let  $z_{LP}$  denote the optimal value of (LP).

Show that  $z_{IP} \ge z_D \ge z_{LP}$ .....(10p)

**5.** Consider the linear program (LP) given by

(LP) minimize 
$$x_1 - x_2 + x_3 + x_4$$
  
subject to  $-x_1 + 2x_2 + x_3 - 3x_4 = 1,$   
 $-1 \le x_j \le 1, \quad j = 1, \dots, 4.$ 

Solve (LP) by Dantzig-Wolfe decomposition. Consider  $-x_1 + 2x_2 + x_3 - 3x_4 = 1$  the complicating constraint. Start with the initial basis corresponding to the extreme points  $(-1 \ 1 \ -1 \ 1)^T$  and  $(-1 \ 1 \ 1 \ -1)^T$ . The subproblems that arise may be solved by any method, that need not be systematic. . . . . . . . . . . . . . . . . (10p)

Good luck!